

# Operations vs. \*-tactics

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## Posets preserving PFA

Thm (König - Y. 2003) —

PFA is preserved under any  $\omega_2$ -closed forcing.

Since then, several broader classes of posets preserving PFA have been found. Such classes are defined in terms of Banach-Mazur games.

## The Banach-Mazur game

For a (separative) poset  $P$  and an ordinal  $\alpha$ ,  $G_\alpha(P)$  denotes the following two-player game: Players choose smaller and smaller  $P$ -conditions in turn. At each limit stage, Player II goes first to choose a condition smaller than the preceding moves. II wins iff she was able to make  $\alpha$  moves without breaking the rule.

$$\begin{array}{ccccccccc} \text{I} & a_0 & a_1 & \dots & a_{\omega+1} & \dots \\ & \geq & \geq & \geq & \geq & \geq & \geq \\ \text{II} & b_0 & b_1 & \dots & b_\omega & b_{\omega+1} & \dots \end{array}$$

In this talk, we only consider the case  $\alpha = \omega_1 + 1$ .

## Strategies, tactics and operations

Def A strategy (for Player II) in  $G_{\text{w+1}}(P)$  is a mapping which suggests a move, knowing the full information on the sequence of preceding moves of Player I (called as "the record"), at each turn of Player II.

I	$a_0$	$a_1$	...	$a_{w+1}$	...	
II	$\tau(a_0)$	$\tau(a_0, a_1)$	$\dots$	$\tau(a_0, \dots, a_w)$	$\tau(a_0, \dots, a_w, \hat{a}_{w+1})$	...

Def A tactic is a strategy whose suggestions depend only on

"the current position" (= the Boolean infimum of "the record").

I	$a_0$	$a_1$	...	$a_{w+1}$	...	
II	$T(a_0)$	$T(a_1)$	...	$T(\bigwedge a_n)$	$T(a_{w+1})$	...

Def An operation is a strategy whose suggestions depend only on

"the current position" and "the number of the turn" ( $\equiv$  the length of the record).

I	$a_0$	$a_1$	...	$a_{w+1}$	...	
II	$T(0, a_0)$	$T(1, a_1)$	...	$T(w, \bigwedge a_n)$	$T(w+1, a_{w+1})$	...

## The modified game $G^*(P)$

In this game, Player I chooses a countable set of P-conditions, instead of a single condition at each turn.

$$\begin{array}{ccccccc} \text{I} & A_0 (< [P]^{<\omega}) & A_1 & \cdots & A_{\omega+1} & \cdots \\ \text{II} & b_0 & b_1 & \cdots & b_\omega & b_{\omega+1} & \cdots \end{array}$$

They must obey the following inequalities:

$$\wedge A_0 \geq b_0 \geq \wedge(A_0 \cup A_1) \geq b_1 \geq \cdots \geq \wedge(\bigcup_{n<\omega} A_n) \geq b_\omega \geq \wedge((\bigcup_{n<\omega} A_n) \cup A_{\omega+1}) \geq b_{\omega+1} \geq \cdots$$

That is, the Boolean infimum of the union of preceding moves of Player I at each moment plays the role of I's move in the usual B-M game.

II wins iff she was able to make w+1 moves without breaking the rule.

## \*-tactics and \*-operations

The notion of a strategy for  $G^*(P)$  is defined in the same way as that of  $G_{w+1}(P)$ .

Def A \*-tactic (for  $G^*(P)$ ) is a strategy whose suggestions depend only on the union of "the record".

$$\begin{array}{ccccccc} \text{I} & A_0 & A_1 & \cdots & A_{w+1} & \cdots \\ \text{II} & T(A_0) & T(A_0 \cup A_1) & \cdots & T(\bigcup_{n \leq w} A_n) & T(\bigcup_{n \leq w} A_n \cup A_{w+1}) & \cdots \end{array}$$

A \*-operation is a strategy whose suggestions depend only on the union of the record and "the number of the turn".

## Game closure properties

Def A poset  $P$  is  $(\omega+1)$ -

} strategically  
tactically  
operationally  
\*-tactically  
\*-operationally

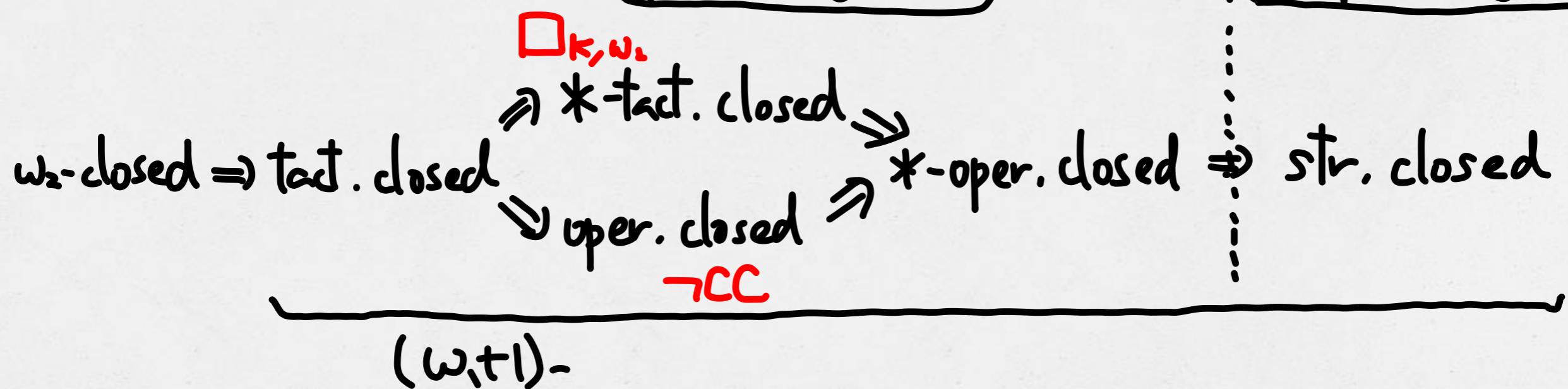
closed if II has a

winning { strategy  
tactic  
operation  
\*-tactic  
\*-operation }  
(for  $G_{\omega+1}(P)$ , or  $G^*(P)$  respectively).

Thm (Y, 2011) PFA is preserved under any  
 $(\omega+1)$ -\*-operationally closed forcing.

preserving PFA

not preserving PFA



Question Compare  $(\omega_1 + 1)$ -oper. closed and  $(\omega_1 + 1)$ -\*-tadt. closed.

## MA<sub>ω\_1</sub>

Def CP  $\Leftrightarrow \exists f: \omega_1 \rightarrow \omega_1, \forall \beta \in \Sigma^2_1 (= \{\beta < \omega_1 \mid cf\} = \omega_1)$

$\exists C \subseteq_{\text{dub}} \beta [0.t.C = \omega_1 \wedge \forall \beta \in C (f(\beta) = 0.t.(C \cap \beta))]$ .

SCP  $\Leftrightarrow \exists \langle z_\alpha \mid \alpha \in \Sigma^2_0 \rangle$  s.t.

-  $\forall \alpha \in \Sigma^2_0 [z_\alpha \subseteq_{\text{cofinal}} \alpha \wedge |z_\alpha| = \omega]$ ,

-  $\forall \beta \in \Sigma^2_1 \exists C \subseteq_{\text{dub}} \beta \cap \Sigma^2_0$  s.t.

0.t. C = ω<sub>1</sub>  $\wedge \langle z_\alpha \mid \alpha \in C \rangle$  is a  $\subseteq$ -continuous increasing sequence.

Fact  $\text{CP} \equiv \text{MA}_{\omega_1}((\omega_1+1)\text{-oper. closed})$ .

$\text{SCP}^- \equiv \text{MA}_{\omega_1}((\omega_1+1)\text{-* -tact. closed})$ .

Remark A natural poset forcing  $\{\begin{matrix} \text{CP} \\ \text{SCP} \end{matrix}\}$  is  
 $(\omega_1+1)\text{-}\left\{\begin{array}{l} \text{oper.} \\ \text{* -tact.} \end{array}\right\}\text{-closed.}$

## Main Results

Thm (Y., 2012) Under  $\text{MA}^+(\omega, \text{-closed})$

- (1)  $\text{IT}_P \supset \text{CP}$  for any  $(\omega+1)$ -\*-tact. closed  $P$ .
- (2)  $\text{IT}_P \supset \text{SCP}^-$  for any  $(\omega+1)$ -oper. closed  $P$ .

Cor Under  $\text{MA}^+(\omega, \text{-closed})$ .

- (1) There exists an  $(\omega+1)$ -oper. closed poset which is not  $(\omega+1)$ -\*-tact. closed.
- (2) There exists an  $(\omega+1)$ -\*-tact. closed poset which is not  $(\omega+1)$ -oper. closed.

## Remark

As for (1) of the main theorem, in fact we can show that under  $\text{MA}^+(\omega_1\text{-closed})$ .

$\Vdash_{\mathbb{P}} \text{CC}$  for any  $(\omega_1+1)\text{-*}\text{-tact. closed }$   $\mathbb{P}$ .

(Note that CC negates CP).

This can be proved by imitating Miyamoto's proof of the consistency of  $\text{CC} + \text{the Strong Non Reflection}$  using  $\text{MA}^+(\omega_1\text{-closed})$ , which was inspired by Sakai's proof of the consistency of  $\Box_{\omega_1, 2}$  and CC.

## Some ideas for proof of (2)

$P$ :  $(\omega+1)$ -oper. closed poset,  $\sigma$ : winning operation for  $G_{\omega+1}(P)$ .

$R$ : the poset of all partial plays (in  $G_{\omega+1}(P)$ ) ending with  $\Pi$ 's move, where  $\Pi$  follows  $\sigma$ .

$R$  has a natural projection to  $P$   
 $(\pi(r))$  is the last move of  $\Pi$  in  $r$  for  $r \in R$  ) and thus  
is forcing equivalent to  $P * \dot{Q}$  for some  $\dot{Q}$ .

It is not hard to see that (2) follows from the following lemma:

Lemma Let  $\langle \dot{z}_\alpha : \alpha \in \mathbb{S}_0^2 \rangle$  be a sequence of IP-names s.t.

$$\Vdash_{\text{IP}} \forall \alpha \in \mathbb{S}_0^2 (\dot{z}_\alpha \subseteq_{\text{cofinal}} \check{\omega} \wedge |\dot{z}_\alpha| = \omega).$$

Then in  $V^{\text{IP}*Q} = V^{\mathbb{R}}$ ,

$S = \{x \in [\omega_1^\text{v}]^\omega \mid \dot{z}_{\sup(x)} \not\subseteq \check{x}\}$  is stationary in  $[\omega_1^\text{v}]^\omega$ .

Note  $\dot{\Theta}$  is not necessarily proper in  $V^{\text{IP}}$ .

Sketch of proof of Lemma Work in  $V$ .

$\theta$ : large regular.  $\mathbb{A} = \langle H_\theta, \in, \dots \rangle$  enough structured.

It is enough to show that for each  $r \in R$ , there exists a countable elementarily substructure  $N$  and an  $(N, R)$ -generic condition  $s \leq r$  s.t.

s.t.  $s \Vdash_{R^+} \dot{z}_{\sup(\check{v}_n w_2^v)} \not\in \check{v}$  ( $\Leftrightarrow \pi(s) \Vdash_P \neg r$ ).

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Sublemma Suppose  $N \prec_{\text{ctble}} R$ ,  $r \in R$  and  $(p_n)_{n \in \omega}$  is an

$(N, P)$ -generic sequence with  $\pi(r) \geq p_0$ . Then for each  $q \in P$

extending  $\sigma(N_n w_1, \bigwedge_{n \in \omega} p_n)$  there exists an  $(N, R)$ -generic  $s \leq r$  s.t.  $\pi(s) \leq q$ .

Using a variation of Velickovic games, we can find

$N_0, N_1 \prec_{\text{ctbl}} \lambda$  and  $\langle p_n^0 | n < \omega \rangle, \langle p_n^1 | n < \omega \rangle$  s.t.

(a)  $\sup(N_0 \cap \omega_2) = \sup(N_1 \cap \omega_2)$  ( $=: \gamma$ ).

(b)  $N_0 \cap \omega_1 = N_1 \cap \omega_1$ .

(c)  $N_0 \cap N_1 \cap \omega_2$  is bounded in  $\gamma$ .

(d)  $\langle p_n^i | n < \omega \rangle$  is an  $(N_i, P)$ -generic sequence (for  $i = 0, 1$ ).

(e)  $\pi(r) \geq p_0^0 \geq p_0^1 \geq p_1^0 \geq p_1^1 \geq p_2^0 \geq p_2^1 \geq \dots$

By (b) and (e) we have  $\sigma(N_{0 \cap w_1}, \bigwedge_n p_n^o) = \sigma(N_{1 \cap w_1}, \bigwedge_n p_n^!) (=: g)$ .

By (a) and (c),  $\mathcal{F} \Vdash_P "z; \notin \check{N}_0 \text{ or } z; \notin \check{N}_1"$ . Thus w.m.o.

there exists  $\mathcal{F}' \leq_P \mathcal{F}$  s.t.  $\mathcal{F}' \Vdash_P "z; \notin \check{N}_0"$ . By sublemma with (d) and the fact that  $g = \sigma(N_{0 \cap w_1}, \bigwedge_n p_n^o)$ , there exists an  $(N_0, R)$ -generic  $s \in_R r$  with  $\pi(s) \leq \mathcal{F}' \Vdash "z; \notin \check{N}_0"$ . //